

A logical approach to Isomorphism Testing and Constraint Satisfaction

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ESLLI 2016, 15–19 August

Part 3: $\text{FO}_{\#}^2$ -definable graphs.

Outline

- 1 Warm-up
- 2 Local structure of $\text{FO}_{\#}^2$ -definable graphs
- 3 Global structure of $\text{FO}_{\#}^2$ -definable graphs
- 4 References

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- Every discrete graph; hence, almost all graphs.

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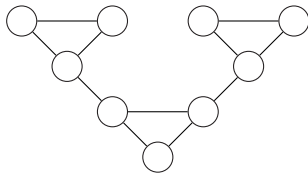
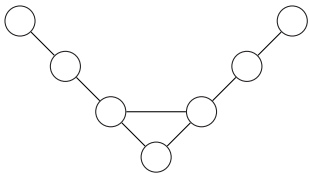
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- Every **unigraph**, i.e., every graph characterizable by its degree sequence up to isomorphism, like K_n , $K_{1,n}$, mK_2 , C_4 , $C_5 \dots$

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- Every discrete graph; hence, almost all graphs.
- Every **unigraph**, i.e., every graph characterizable by its degree sequence up to isomorphism, like K_n , $K_{1,n}$, mK_2 , C_4 , $C_5 \dots$
- Every tree; because every tree is characterizable by Color Refinement [Edmonds 65] (a proof is postponed).

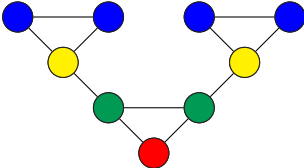
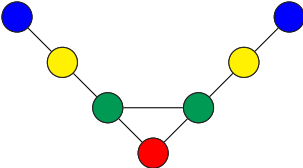
Example

Which one of the following two graphs is definable in $\text{FO}_{\#}^2$?



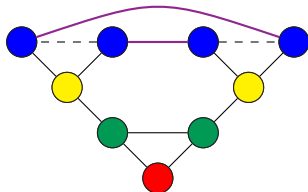
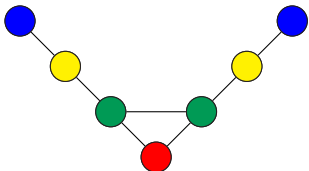
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Local structure of definable graphs

Definition

Call the stable color classes of vertices **cells**.

Let X be a cell in G .

- $G[X]$ is regular.
- If G is definable, then $G[X]$ must be a unigraph (to be proved shortly).

Complete list of regular unigraphs [Johnson 1975]

- Complete and empty graphs. (homogeneous)
- Matchings mK_2 and their complements. (heterogeneous)
- The cycle C_5 . (heterogeneous)

Local structure of definable graphs

Let X, Y be cells in G .

- $G[X, Y]$ is bi-regular.
- If G is definable, then $G[X, Y]$ is identified by its degrees.

Complete list of such bi-regular graphs [Koren 1976]

- Complete bipartite graphs and empty graphs.
(isotropic)
- Forests of stars $sK_{1,t}$, $s \geq 2$, and their bipartite complements.
(anisotropic)

Local structure of definable graphs

Lemma

If G is definable, then for any cells X, Y

- (A) $G[X]$ is from the Johnson list.
- (B) $G[X, Y]$ is from the Koren list.

Proof of (A)

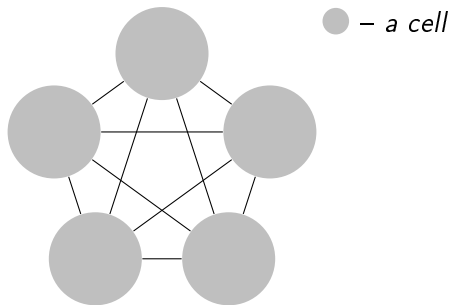
Assume that $G[X]$ is not a unigraph, that is, there is a regular graph $F \not\cong G[X]$ of the same degree with as many vertices. Change G on X so that $G'[X] \cong F$, where G' denotes the modified graph. Then

- For every vertex v , $C^i(v)$ is the same in G and G' (induction on i);
- Therefore, CR does not distinguish G and G' ;
- $G \not\cong G'$ because any isomorphism respects the color classes and should be an isomorphism also between $G[X]$ and $G'[X]$.

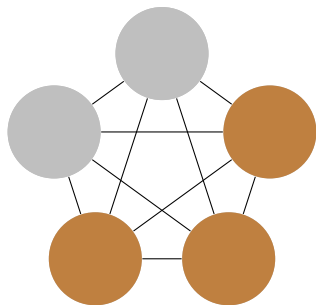
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Cell-Graph

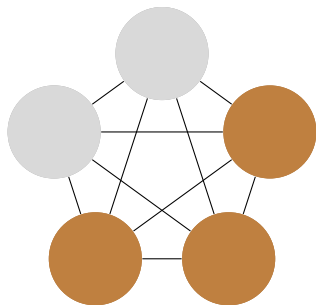


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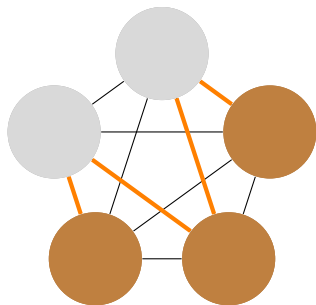
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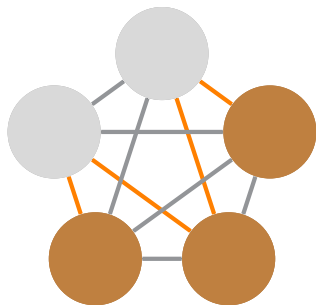


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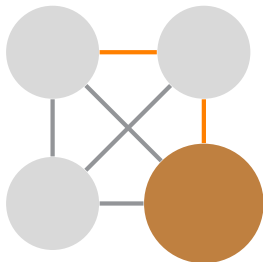
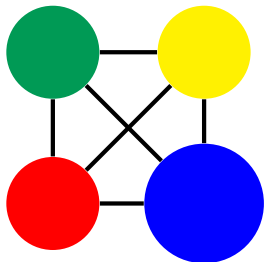
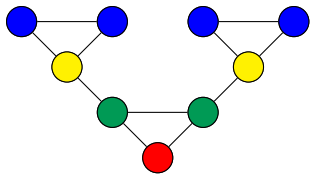


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- Homogeneous cells (empty/complete),
- Anisotropic edges ($sK_{1,t}$ and bipartite complements)
- Isotropic edges (empty/complete bipartite)

Example



Global structure of definable graphs

Configurations forbidden in the cell-graph of a definable graph:

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(C) An anisotropic path of equally sized cells connecting two heterogeneous cells:



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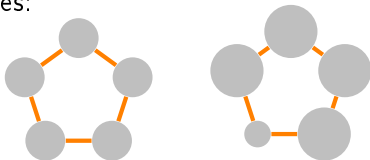
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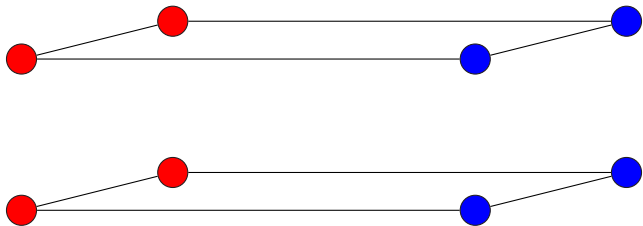
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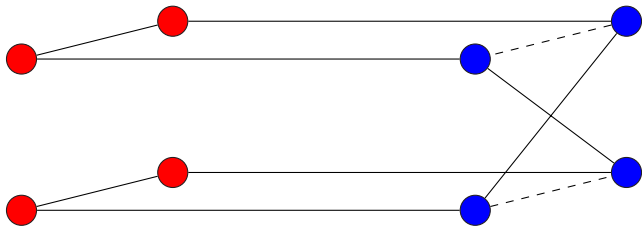
(F) Anisotropic cycles:



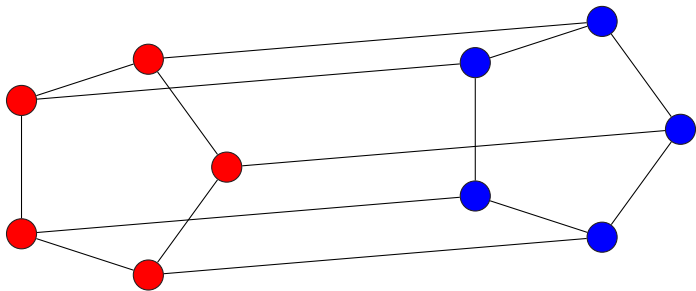
Proof-sketch of Condition (C)



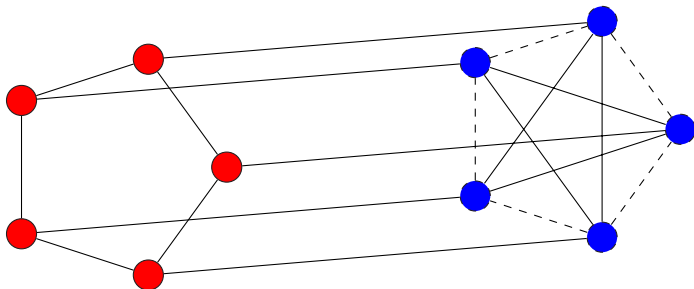
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Proof-sketch of Condition (C)



Characterization of definable graphs

Theorem (Arvind, Köbler, Rattan, V. 2015 and Kiefer, Schweitzer, Selman 2015)

Conditions (A)–(F) are both necessary and sufficient for definability of a graph in $\text{FO}_{\#}^2$.

Corollary

The class of graphs definable in $\text{FO}_{\#}^2$ is recognizable in polynomial time.

Exercise

Using the characterization of $\text{FO}_{\#}^2$ -definable graphs, prove that every forest is definable in $\text{FO}_{\#}^2$.

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- V. Arvind, J. Köbler, G. Rattan, and O. Verbitsky. On the power of color refinement. FCT'15.
- S. Kiefer, P. Schweitzer, and E. Selman. Graphs identified by logics with counting. MFCS'15.